

## Equivalence Relations and Partitions:

↑  
generalize the idea of equality

An equivalence relation on a set  $X$  is a relation  $R \subseteq X \times X$  s.t.

- $(x, x) \in R \quad \forall x \in X$  (reflexive Property)
- $(x, y) \in R \Rightarrow (y, x) \in R$  (symmetric Property)
- $(x, y) \in R$  and  $(y, z) \in R \Rightarrow (x, z) \in R$   
(transitive Property)

Often write  $(x, y) \in R \Leftrightarrow x \sim y$

- $x \sim x$
- $x \sim y \Rightarrow y \sim x$
- $x \sim y$  and  $y \sim z \Rightarrow x \sim z$

A partition  $\mathcal{P}$  of a set  $X$  is a collection of non-empty sets  $X_1, X_2, \dots$  s.t.  $X_i \cap X_j = \emptyset$  for  $i \neq j$   
and  $\bigcup_k X_k = X$

### Equivalence class

$$[x] = \{ y \in X : y \sim x \}$$

↑  
 $x$  is the representative of the eq. class.

**Theorem 1.25.** Given an equivalence relation  $\sim$  on a set  $X$ , the equivalence classes of  $X$  form a partition of  $X$ . Conversely, if  $\mathcal{P} = \{X_i\}$  is a partition of a set  $X$ , then there is an equivalence relation on  $X$  with equivalence classes  $X_i$ .

Proof:

First suppose  $\exists$  eq. relation  $\sim$  on  $X$

$\Rightarrow$  for any  $x \in X$  we have that  $x \in [x]$  (by the reflexive property)

$\therefore [x]$  is non-empty

$$X = \bigcup_{x \in X} [x]$$

Given  $x, y \in X$  it remains to show that either

$$[x] = [y] \text{ or } [x] \cap [y] = \emptyset$$

Suppose  $[x] \cap [y]$  is nonempty  $\Rightarrow \exists z \in [x] \cap [y]$

$z \sim x$  and  $z \sim y$  by symmetry and transitivity

$$x \sim y \Rightarrow [x] \subseteq [y]$$

$$\text{and } y \sim x \Rightarrow [y] \subseteq [x]$$

$$[x] = [y]$$

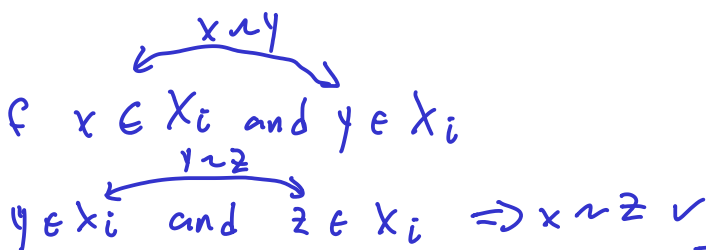
$\therefore$  Two eq. classes are either disjoint or the same

Now suppose that  $\mathcal{P} = \{X_i\}$  is a partition

Let  $x \sim y$  if  $x \in X_i$  and  $y \in X_i$

$$x \sim x \quad \checkmark$$

If  $x \sim y \Rightarrow y \sim x$  by def. If  $x \in X_i$  and  $y \in X_i$



Corr 1 Two eq. classes are either disjoint or equal

Ex 1 Let  $p, q, r, s \in \mathbb{Z}$ ,  $q \neq 0, s \neq 0$

Define  $\frac{p}{q} \sim \frac{r}{s}$  if  $ps = qr$  (i.e.  $\frac{1}{2} \sim \frac{2}{4}$ )

$\frac{p}{q} \sim \frac{p}{q} \therefore$  reflexive  $\frac{p}{q} \sim \frac{r}{s} \Rightarrow \frac{r}{s} \sim \frac{p}{q} \therefore$  symmetric

Suppose  $\frac{p}{q} \sim \frac{r}{s}$  and  $\frac{r}{s} \sim \frac{t}{u}$  ( $q, s, u \neq 0$ )

Then  $ps = qr$  and  $ru = st$   
 $\Rightarrow$   $psu = qru = qst$

$psu = qst \Rightarrow pu = qt \Rightarrow \frac{p}{q} \sim \frac{t}{u}$

$(p, q)$  and  $(r, s)$  are in the same class if they reduce to the same fraction in lowest terms

Ex 2. Let  $r, s$  be in  $\mathbb{Z}$ , suppose  $n \in \mathbb{N}$

$r$  is congruent to  $s$  modulo  $n$

or

$r \equiv s \pmod{n}$  ( $r = s \pmod{n}$ )

If  $r - s = n \cdot k$  for some  $k \in \mathbb{Z}$

(alt.  $r - s$  is evenly divisible by  $n$ )

$$\text{i.e. } 41 \equiv 17 \pmod{8}$$

$$\text{Since } 41 - 17 = 24 = 8 \cdot 3$$

Congruence mod  $n$  forms an Eq. relation on  $\mathbb{Z}$

$$\bullet \quad r \equiv r \pmod{n} \quad \text{since } n - r = 0 = 0 \cdot n$$

$$\bullet \quad r \equiv s \pmod{n} \Rightarrow r - s = n \cdot k \Rightarrow s - r = n \cdot (-k) \\ s \equiv r \pmod{n}$$

$$\bullet \quad r \equiv s \pmod{n} \text{ and } s \equiv t \pmod{n}$$

$$\Rightarrow r - s = kn \text{ and } s - t = ln \quad (k, l \in \mathbb{Z})$$

$$r - s + s - t = kn + ln$$

$$r - t = n(k+l) \Rightarrow r \equiv t \pmod{n} \quad \therefore \text{transitive.}$$

Consider  $\mathbb{Z}/3\mathbb{Z} \cong$  integers modulo 3, we have

$$[0] = \{ \dots, -3, 0, 3, 6, \dots \}$$

$$[1] = \{ \dots, -2, 1, 4, \dots \}$$

$$[2] = \{ \dots, -1, 2, 5, 8, \dots \}$$

# Integers + Mathematical Induction

**Principle 2.1** (First Principle of Mathematical Induction). Let  $S(n)$  be a statement about integers for  $n \in \mathbb{N}$  and suppose  $S(n_0)$  is true for some integer  $n_0$ . If for all integers  $k$  with  $k \geq n_0$ ,  $S(k)$  implies that  $S(k+1)$  is true, then  $S(n)$  is true for all integers  $n$  greater than or equal to  $n_0$ .

How do we show that

$$(*) \quad 1 + 2 + \dots + n = \frac{n(n+1)}{2} \quad \text{for any } n?$$

$n=1$

$$1 = \frac{1 \cdot (1+1)}{2} \quad \checkmark \text{ True } n=1$$

By Induction we may assume this is true for  $n$  and show that it holds for  $n+1$

Show: (\*) holds for  $n+1$

$$\begin{aligned} \underbrace{1 + 2 + \dots + n}_{= \frac{n(n+1)}{2}} + (n+1) &= \frac{n(n+1)}{2} + n+1 \\ &\simeq \frac{n^2 + 3n + 2}{2} = \frac{(n+1)(n+1+1)}{2} \end{aligned}$$

**Principle 2.5** (Second Principle of Mathematical Induction). Let  $S(n)$  be a statement about integers for  $n \in \mathbb{N}$  and suppose  $S(n_0)$  is true for some integer  $n_0$ . If  $S(n_0), S(n_0+1), \dots, S(k)$  imply that  $S(k+1)$  for  $k \geq n_0$ , then the statement  $S(n)$  is true for all integers  $n \geq n_0$ .

A non-empty subset  $S$  of  $\mathbb{Z}$  is well-ordered if  $S$  contains a least element

$\left( \begin{array}{l} \mathbb{Z} \text{ not well-ordered since} \\ \quad \text{no least element} \\ \mathbb{N} \text{ is well ordered} \end{array} \right)$

**Principle 2.6** (Principle of Well-Ordering). *Every nonempty subset of the natural numbers is well-ordered.*

The Principle of Well-Ordering is equivalent to the Principle of Mathematical Induction.

**Lemma 2.7.** *The Principle of Mathematical Induction implies that 1 is the least positive natural number.*

Proof:

$$S = \{ n \in \mathbb{N} \mid n \geq 1 \} \Rightarrow 1 \in S \quad | \geq 1$$

$$\text{assume } n \in S \quad n \geq 1$$

$$n+1 \geq 1 \Rightarrow n+1 \in S \quad \therefore \text{by induction all}$$

Natural numbers are  $\geq 1$ .

**Theorem 2.8.** *The Principle of Mathematical Induction implies the Principle of Well-Ordering. That is, every nonempty subset of  $\mathbb{N}$  contains a least element.*